

Section 10. Data Storages

- Three Basic Functions of Any Digital System:
 - Data Processing
 - ❖ CPU/Memory/IO process data
 - ❖ Basic computations, searching, sorting, etc
 - Data Transmission
 - ❖ Buses, parallel/serial interfaces
 - ❖ Wired/wireless networking
 - Data Storage
 - ❖ Keeping persistent data
 - ❖ Disks, Flash memory, Tapes etc.



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Data Explosion!



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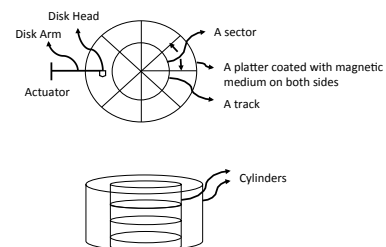
Important Issues in Data Storage

- Performance
 - Improve I/O performance not to drag down CPU and memory
 - ✓ Hard disk drive technology
 - ✓ Disk arrays: RAID architecture
- Reliability
 - Data are important assets of businesses and organizations
 - Keeping data reliable is important
- Recoverability
 - Be able to recover data in case of failure
 - Existing data recovery technology and a new one



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Terminologies in a Hard Disk Drive



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Necessary Steps of a Disk Access

1. **Controller command interpretation**
 - Take and interpret commands, LBA, Sizes from host
2. **Seek**
 - Move disk head to the track where data is stored
3. **Rotation**
 - Wait for the accessed sector to rotate right under the disk head
4. **Data Transfer**
 - Transfer data from disk to controller buffer then to host



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RAID Architectures

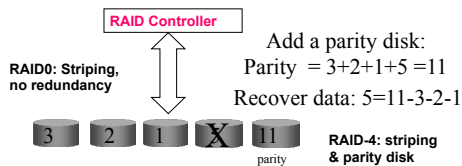
1. **Motivations:**
 - Cost of disks is low, is getting less expensive
 - Improve performance, throughput, by parallelism
 - ✓ Multiple I/O requests are done in parallel
 - ✓ Striping single request into chunks
 - Improve reliability using redundancy
2. **Different RAID Architectures**
 - RAID0: striping, no redundancy
 - RAID1: Mirroring, or shadowing, no striping, 2xN redundancy
 - RAID2: Memory like ECC, Hamming code
 - RAID3: Knowing which disk failed, erasure code, Bit-interleaved,
 - RAID4: Block-interleaved, one dedicated parity disk, striping unit varies
 - RAID5: Block-interleaved, distributed parity, load balancing.



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RAID Architectures

RAID, redundant array of inexpensive disks
For high speed and high reliability

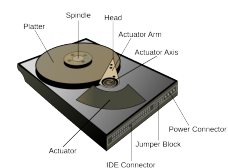
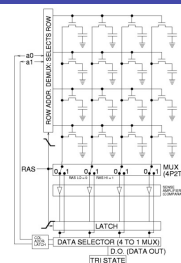


If parity blocks for different stripes are distributed evenly across disks, we got RAID-5



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Traditional Storage Technologies

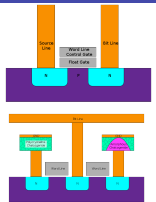


Controller command interpretation: Take and interpret commands, LBA, Sizes from host
Seek: Move disk head to the track where data is stored
Rotation: Wait for the accessed sector to rotate right under the disk head
Data Transfer: Transfer data from disk to controller buffer then to host

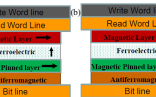


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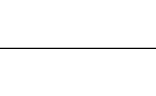
Emerging Storage Technologies



Flash Memory: Store a bit by trapping charges
Read is fast, write is slow and wearing



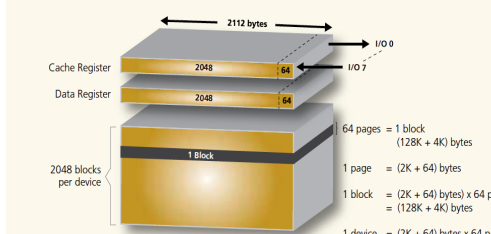
Phase Change Memory:
crystalline and amorphous states of chalcogenide glass have dramatically different resistivity
Much faster switching
Fast read, slow write, wearing, too, but less than flash



Magnetoresistance RAM: fast switching, symmetric, no durability issue
Promising technology once cost comes down

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A typical NAND Flash



Cache Register: 2048
Data Register: 2048
2048 blocks per device

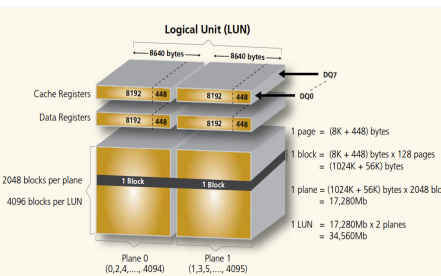
2112 bytes
I/O 0
I/O 7

64 pages = 1 block (128K + 4K) bytes
1 page = (2K + 64) bytes
1 block = (2K + 64) bytes x 64 p = (128K + 4K) bytes
1 device = (2K + 64) bytes x 64 p x 2048 blocks = 2112 Mb

3: 2Gb Array Organization of 90nm Micron SLC NAND Flash (2006)

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Newer SLC NAND Flash



Logical Unit (LUN)
8640 bytes
Cache Registers: 8192, 448
Data Registers: 8192, 448
2048 blocks per plane
4096 blocks per LUN

Plane 0 (0,2,4,..., 4094)
Plane 1 (1,3,5,..., 4095)

1 page = (8K + 448) bytes
1 block = (8K + 448) bytes x 128 pages = (1024K + 56K) bytes
1 plane = (1024K + 56K) bytes x 2048 blocks = 17,280Mb
1 LUN = 17,280Mb x 2 planes = 34,560Mb

Figure 4: 32Gb Array Organization of 25nm Micron SLC NAND Flash (2010)

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